

### Solution #1:

1. Show that  $\Omega(g(n)) \cap o(g(n))$  is the empty set.

Solution: Suppose  $f(n) = o(g(n))$ . This implies that  $\lim_{n \rightarrow \infty} \left[ \frac{f(n)}{g(n)} \right] = 0$ . If  $f(n) = \Omega(g(n))$ , then the previous limit can not be equal to 0. Hence the result follows.

2. Does the statement  $[f(n) = O(g(n))]$  imply the statement  $[2^{f(n)} = O(2^{g(n)})]$ ? Is the converse true? Give proofs or counter examples.

Solution: Let  $f(n) = \lg[n^n] = n \lg n$ ; and  $g(n) = \lg(n!)$ . We have shown that  $f(n) = \Theta(g(n))$ . But we have also shown that  $n^n = \omega(n!)$  and  $\{\omega, O\}$  are incompatible. So  $2^{f(n)} = n^n \neq O(n!)$ ;  $n! = 2^{g(n)}$ . So the implication is not true. The converse would be:  $[2^{f(n)} = O(2^{g(n)})] \implies [f(n) = O(g(n))]$  and this is true as shown below:

If  $f(n) = O(g(n))$ , then we show that  $\lg(f(n)) = O(\lg(g(n)))$

$f(n) \leq c_2 g(n) \implies \lg[f(n)] \leq \lg[c_2 g(n)] = \lg[g(n)] + \lg c_2 \leq 2 \lg[g(n)]$  for large  $n$ .

3. Exercise 3.1-4:

Solution:

(i) Is  $2^{n+1} = O(2^n)$ :  $\lim_{n \rightarrow \infty} \left[ \frac{2^{n+1}}{2^n} \right] = 2$  and therefore the answer is yes.

(ii) Is  $2^{2^n} = O(2^n)$ :  $\lim_{n \rightarrow \infty} \left[ \frac{2^{2^n}}{2^n} \right] = \infty$  and therefore the answer is no

4. Exercises 3.2-4: Are functions  $[\lg n]!$  and  $[\lg \lg n]!$  polynomially bounded?

Solution:

A function  $f(n)$  is said to be polynomially bounded iff  $f(n) = O(n^b)$  for some fixed  $b$ .

Let  $[\lg n] = k$ . This is equivalent to  $[k-1 < \lg n \leq k] \iff 2^{k-1} < n \leq 2^k$ .  
 $[\lg n]! = k!$

$n^b \leq (2^k)^b = 2^{bk} < k!$  since  $bk < k \lg k$  for large values of  $k$ . Hence, the first function is not polynomially bounded.

Let  $[\lg \lg n] = k \iff k-1 < \lg \lg n \leq k \iff 2^{k-1} < \lg n \leq 2^k \iff 2^{2^{k-1}} < n \leq 2^{2^k}$

$[\lg \lg n]! = k!$

$n^b \leq (2^{2^k})^b = 2^{b2^k} > k!$  since  $b2^k > k \lg k$  for large  $k$ . Hence, the second function is polynomially bounded.

5. Show that (i)  $\sum_{i=1}^n i^2 = \Theta(n^3)$

Solution: We start with the relation:

$$\begin{aligned}\sum_{i=1}^{n+1} i^3 &= \sum_{i=0}^n (i+1)^3 \\ &= \sum_{i=0}^n (i^3 + 3i^2 + 3i + 1) \\ &= \sum_{i=0}^{n-1} i^3 + 3 \sum_{i=0}^{n-1} i^2 + 3 \sum_{i=0}^{n-1} i + \sum_{i=0}^{n-1} 1\end{aligned}$$

Hence

$$\sum_{i=0}^{n-1} i^2 = \frac{2n^3 - 3n^2 + n}{6}$$

Hence

$$\begin{aligned}\sum_{i=1}^n i^2 &= \frac{2(n+1)^3 - 3(n+1)^2 + n + 1}{6} \\ &= \frac{2n^3 + 3n^2 + n}{6} \\ &= \frac{n(n+1)(2n+1)}{6} \\ &= \Theta(n^3)\end{aligned}$$

This gives you the method of finding  $\sum_{i=1}^n i^k$  for general values of  $k$  by induction.