

1. Let  $x^0$  be an optimal solution to the linear program:

$$\begin{aligned} & \min cx \\ Ax & = b; x \geq 0 \end{aligned}$$

and let  $y^0$  be an optimal solution to its dual. Now consider the LP which is the same as above except the  $i^{th}$  constraint in the new problem is the old constraint multiplied by  $k \neq 0$ . Can you find an optimal solution to the new problem and its dual in terms of the solutions to the old primal and dual and the value of  $k$ ?

2. Solve M.Beale's example shown below by lexicographic method:

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$-z$	$=$	rhs
$\frac{1}{4}$	-60	$-\frac{1}{25}$	9	1	0	0	0	=	0
$\frac{1}{2}$	-90	$-\frac{1}{50}$	3	0	1	0	0	=	0
0	0	1	0	0	0	1	0	=	1
$-\frac{3}{4}$	150	$-\frac{1}{50}$	6	0	0	0	1	=	0

This example cycles if regular simplex method is used.

3. Consider the following problem known in the literature as the two-person-zero-sum game:

$$\begin{aligned} & \min M \\ & \sum_{j=1}^n a_{i,j}y_j \leq M; 1 \leq i \leq m \\ & \sum_{j=1}^n y_j = 1; y_j \geq 0; 1 \leq j \leq n \end{aligned}$$

Show that its dual is:

$$\begin{aligned} & \max N \\ & \sum_{i=1}^m a_{i,j}x_i \geq N; 1 \leq j \leq n \\ & \sum_{i=1}^m x_i = 1; x_i \geq 0; 1 \leq i \leq m \end{aligned}$$

Prove that :

- (i)  $N \leq M$  for any pair of feasible solutions. .
  - (ii) both problems are feasible and hence have optimal solutions and  $N^* = M^*$ .
  - (iii)  $N^* = M^* > 0$  if  $A > 0$ .
4. Prove that the optimal dual solution is never unique if the optimal primal basic solution is degenerate and the optimal dual is not.

5. Let the LP:  $[\min cx : Ax \geq b; x \geq 0]$  have an optimal solution. Let the LP with the same constraint matrix and objective function but with  $b'$  as the right hand side be feasible. Show that it has an optimal solution if it is feasible.